Effective height bounds for odd-degree totally real points on some curves.

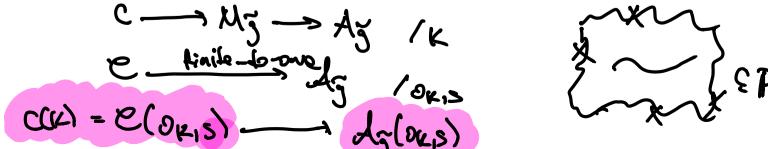
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Diophantus, may your soul be with Satan because of the difficulty of the other theorems of yours, and in particular of the present theorem. — Chortasmenos, ~ 1400.

Book II of Arithmetica:  
solve 
$$y^2 = x^6 + x^2 + 1$$
 non-trivially.  
finds  $x = \frac{1}{2}$ .  
(98 PhD thesis of Wetherell: Drophantus found ble only  
nontrivi. solves.

## Theorem (Faltings). Let $K/\mathbb{Q}$ be a number field. Let C/K be a smooth projective hyperbolic curve. Then: C(K) is finite.

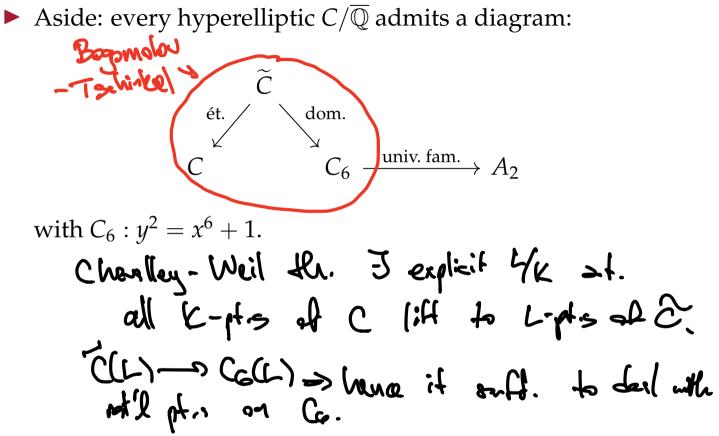
Famously ineffective, i.e. does *not* provide a finite-time algorithm to find C(K) given K and C/K.





Faltings' starting point was Parshin's key observation: for all hyperbolic  $C/\overline{\mathbb{Q}}$ , there is a  $g \ge 2$  and a finite-to-one map  $C \to A_g$ , thanks to an amazing construction of Kodaira.

But nothing forces us to use these maps...



► Trick:  $C_6/\overline{\mathbb{Q}}$  is a Shimura curve, corresponding to  $[G,G] \subseteq G$  with  $G := \Delta(2,6,6)$ , an arithmetic triangle group. (Quat. alg.:  $B_6/\mathbb{Q}$ , indef. of disc. 6.)

Key point of this talk: introducing modularity into Faltings' technique.

Currently, modularity is best understood for GL<sub>2</sub>. Thus we will map to Hilbert modular varieties instead of A<sub>g</sub>.

Downside: unclear for which curves one can (or can't) do invite: this!  $F = Q(G_2),$   $O = O_F = Z_2G_2$   $H_O \subseteq \mathbb{R}^N,$  dim.  $H_O = \mathcal{I}_F:Q_2.$  $dim. \partial H_O = O.$ 

Let  $F/\mathbb{Q}$  be a totally real field. Let  $\mathfrak{o} \subseteq F$  be an order. Let  $K/\mathbb{Q}$  be totally real with  $[K : \mathbb{Q}]$  odd. Let S be a finite set of places of K. Then: there is an effectively computable  $h_{\mathfrak{o},K,S} \in \mathbb{Z}^+$  such that, for all  $[F : \mathbb{Q}]$ -dimensional abelian varieties A/K with good reduction outside S and admitting  $\mathfrak{o} \hookrightarrow \operatorname{End}_K(A)$ ,

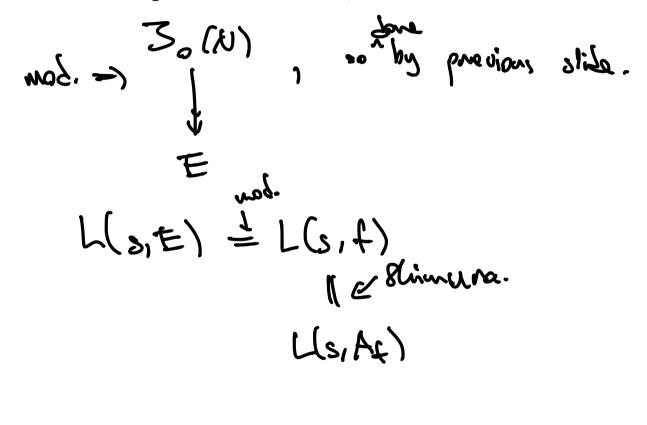
 $h(A) \leq h_{\mathfrak{o},K,S}.$ 

▶ Thus all  $P \in \mathcal{H}_{\mathfrak{o}}(\mathfrak{o}_{K,S})$  satisfy  $h(P) \leq h_{\mathfrak{o},K,S}$ .

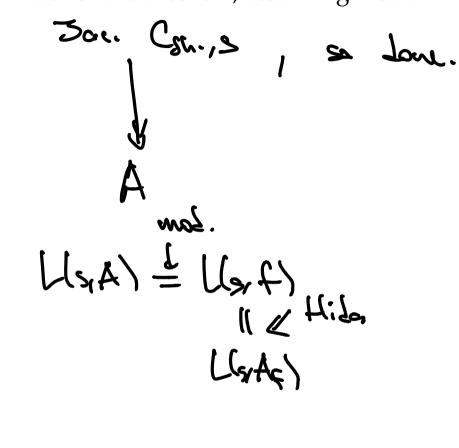
►  $K = \mathbb{Q}$  and  $\mathfrak{o} = \mathbb{Z}$ : Murty-Pasten. Independently,  $K = \mathbb{Q}$  and all  $\mathfrak{o}$ : von Känel.

Masser-Wüstholz/Bost" argument:  
given 
$$B/k$$
, wond all  $k$ -iog. factors  $A/k$ .  
Poinceni =  $3C/k$  e.f.  $B \sim k A \times C$ .  
 $faugre/M.-W. = h(A \times C) \ll 1$ .  
 $I$   
 $Boot = h(A) + h(C)$   
 $W(A) - (0^{10} \cdot Jim. B)$   
 $= h(A) \notin 1$ .  
 $B$ 

Example: let's find all  $E/\mathbb{Q}$  of conductor N.



Proof of the theorem, assuming modularity:



So: done if  $\exists$  explicit finite  $\mathcal{F}$  of odd-deg. tot. real L/K s.t. every such ab. var. A/K is modular over some  $L \in \mathcal{F}$ .

Every such A/K is *potentially* modular (Taylor, we use Snowden).

► Roughly: let  $\lambda$  s.t.  $\overline{\rho} := \overline{\rho}_{A,\lambda} : \text{Gal}(\overline{\mathbb{Q}}/K) \to \text{GL}_2(\mathfrak{o}/\lambda)$  has large image. Then there is  $L_{\overline{\rho}}/K$  tot. real of odd deg. s.t.  $A/L_{\overline{\rho}}$  modular.

 $L_{\overline{\rho}}/K$  determined following Taylor's proof. Key step: producing a tot. real point on a certain variety using Moret-Bailly/Rumely (— can just brute force!!). • (~ Lem. of Dimitrov:) once Nm  $\lambda \gg_{\mathfrak{o},K,S} 1$  (explicit),  $\overline{\rho}$  has large image.

Hermite-Minkowski: the possible p
 are explicitly determined (p
 gives a num. field of bdd. deg. that is unram. outside S and Nm λ).

► Thus by doing the brute force (or cleverer) for each such p̄ we produce *F*, QED.

Example: let K/Q tot. real of odd deg.,  $a \in K^{\times}$ ,  $C_a: x^6 + 4y^3 = a^2.$   $x^6 + 4y^3 = 4y^3 = 4y^3.$ 

► Hypergeom. fam. assoc. to  $\Delta(3, 6, 6)$  (arithmetic!): let  $f: C_a \to \mathbb{P}^1$  via  $f(P) := \frac{x(P)^6}{a^2}$ , and let  $A_P$  be the 2-dim.'l quot. of Jac. of (desing. of)  $y^6 = x^4(1-x)^3(1-f(P) \cdot x)$ .

▶ Because the eq. is  $y^6 = \cdots$ , we see  $\mathbb{Z}[\zeta_3] \hookrightarrow \operatorname{End}_{K(\zeta_3)}(A_P)$ . So  $A_P$  is  $\operatorname{GL}_2(\mathbb{Z}[\zeta_3])$ -type over  $K(\zeta_3)$ . Not enough...

$$x^{6} + ctp^{3} = 1$$

 $y^{e} = x^{t} (1 - x)^{2} (1 - f.x)$ dim Ap = 2, p: Sd. (TVKg) >> GSpy (O)  $e \cong e' \oplus c'',$ Lin.p'= Lin. pri= 2.  $F_1(t) = \eta(t) \cdot \eta(t)$ ,  $\eta$  order 6 class.

► Arith. of  $\Delta(3, 6, 6)$  means  $B_6 \hookrightarrow \operatorname{End}_{\overline{\mathbb{Q}}}^0(A_P)$ , so there are other endo.s... In fact (compare Frob. traces of the 2-dim'l rep.s over  $K(\zeta_3)$ )  $A_P/K$  is GL<sub>2</sub>-type over K.

Easy to compute an explicit  $\mathcal{F}$  s.t. each  $A_P$  is  $GL_2(F)$ -type over K for some  $F \in \mathcal{F}$ , done.

## Thanks!