

Proposition

Let $\mathcal{C} = \{G_n : n \in \mathbb{N}\}$ be a class of finite graphs such that:

- (a) Each graph G_n is r -regular (resp. d_n -regular)
- (b) $\text{girth}(G_n) \rightarrow \infty$

Then, every infinite ultraproduct M of graphs in \mathcal{C} is a model of T_r (resp. a model of T_∞ if $d_n \rightarrow \infty$).

Theorem (G., Robles)

Such classes of finite graphs exists, therefore the theories T_r and T_∞ are both pseudofinite.

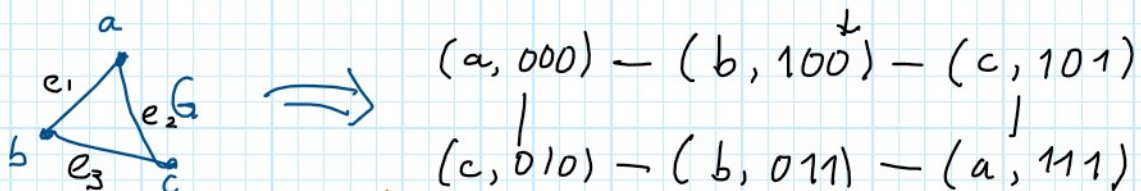
Lifting: $G \xrightarrow{\quad} L[G]$

$G = (V, E)$

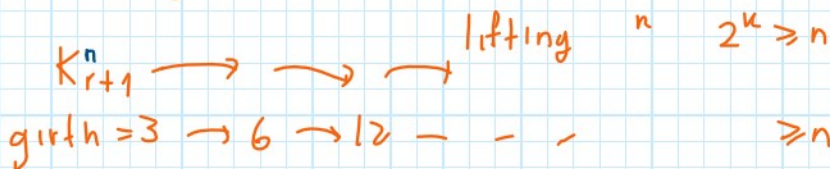
Vertices of $L[G]$: $V \times \{0,1\}^{E(G)}$ $(v, f : E(G) \rightarrow \{0,1\})$

$(v, f) \sim (w, g) \iff e = \{v, w\} \in E \text{ \& } f|_{\neq e} = g|_{\neq e} \text{ \& } f(e) \neq g(e)$

- If G is d -regular $\Rightarrow L[G]$ d -regular
- If $\text{girth}(G) = k \Rightarrow \text{girth}(L[G]) = 2k$



$G_n \xrightarrow{\text{r-regular}} \text{girth}(G_n) \geq n$



Theorem (G., Robles)

Let $\mathcal{C} = \{G_n : n \in \mathbb{N}\}$ be a class of finite graphs such that each graph G_n is d_n -regular and $d_n, \text{girth}(G_n) \rightarrow \infty$.

Let M be an infinite ultraproduct of graphs in \mathcal{C} (a model of T_∞) and fix the non-standard integers $\alpha = |M|$ and $\beta = [d_n]_{\mathcal{U}}$. Then for every formula $\phi(\bar{x}, \bar{y})$ in the language of graphs there is a finite number of polynomials $p_1(X, Y), \dots, p_k(X, Y) \in \mathbb{Z}[X, Y]$ such that:

- ① For every $\bar{a} \in M^{|\bar{x}|}$, $|\phi(M^{|\bar{x}|}, \bar{a})| = p_i(\alpha, \beta)$ for some $i \leq k$.
- ② Moreover, there are formulas $\psi_1(\bar{y}), \dots, \psi_k(\bar{y})$ such that for every $\bar{a} \in M^{|\bar{y}|}$,

$$M \models \psi_i(\bar{a}) \Leftrightarrow |\phi(M^{|\bar{x}|}, \bar{a})| = p_i(\alpha, \beta).$$

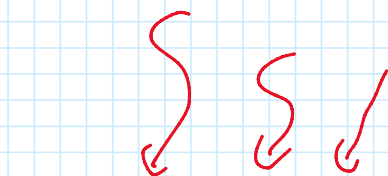
This is enough to show that any class of graphs with the properties above is a multidimensional exact class.

1. Enough to check ^{formulas} one-variable

2. $\text{QE} \rightarrow \Psi(x, \bar{y}) = \bigvee_j \underbrace{\bigwedge_i D_{k_i}(x, y_i)^{n_i}}_{\text{disjunctions can be avoided}}$

3. disjunctions can be avoided

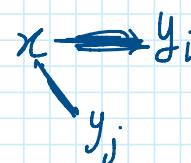
4. $\neg \underline{D}_k(x, a) \wedge \underline{D}_m(x, b) \Leftrightarrow \neg \left(\bigvee D_m(x, b) \wedge D_j(x, a) \right)$



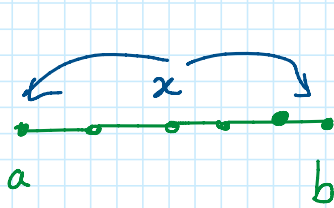
$$D_{k_i}(x, y_i)$$

"possible configurations of y_i 's"

"possibilities for"



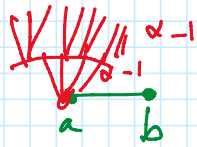
$$D_2(x, a) \wedge D_3(x, b)$$



Polynomial $\psi(y)$

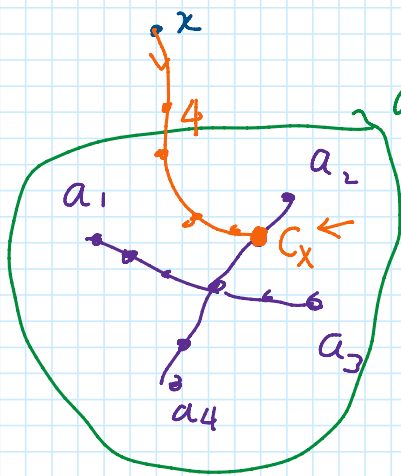
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$$D_5(a, b)$$



$$\underbrace{(d-1)^2}_{D} \quad D_1(a, b)$$

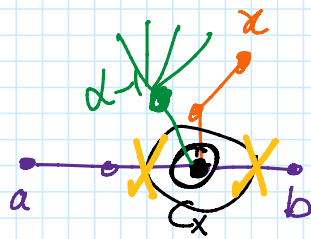
D



$acl(a_1, \dots, a_4)$

$$D_{k_1}(x, a_1) \wedge \dots \wedge D_{k_4}(x, a_4)$$

$$(d-1)^3 \cdot (d - \text{paths in } acl)$$



$$D_4(x, a) \wedge D_3(x, b)$$

$$(d-1)(d-2)$$